Minimum Eccentric Connectivity Index for Graphs with Fixed Order and Fixed Number of Pending Vertices

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Abstract

The eccentric connectivity index of a connected graph $G$ is the sum over all vertices $v$ of the product $d_G(v)e_G(v)$, where $d_G(v)$ is the degree of $v$ in $G$ and $e_G(v)$ is the maximum distance between $v$ and any other vertex of $G$. This index is helpful for the prediction of biological activities of diverse nature, a molecule being modeled as a graph where atoms are represented by vertices and chemical bonds by edges. We characterize those graphs which have the smallest eccentric connectivity index among all connected graphs of a given order $n$. Also, given two integers $n$ and $p$ with $p \leq n - 1$, we characterize those graphs which have the smallest eccentric connectivity index among all connected graphs of order $n$ with $p$ pending vertices.

1 Introduction

A chemical graph is a representation of the structural formula of a chemical compound in terms of graph theory where atoms are represented by vertices and chemical bonds by edges. Arthur Cayley [1] was probably the first to publish results that consider chemical graphs. In an attempt to analyze the chemical properties of alkanes, Wiener [11] has introduced the path number index, nowadays called Wiener index, which is defined as the sum of the lengths of the shortest paths between all pairs of vertices. Mathematical properties and chemical applications of this distance-based index have been widely researched.

Numerous other topological indices are used for quantitative structure-property relationship (QSPR) and quantitative structure-activity relationship (QSAR) studies that help to describe and understand the structure of molecules [3,10], among which the eccentric connectivity index which can be defined as follows. Let $G = (V,E)$ be a simple connected undirected graph. The distance $dist_G(u,v)$ between two vertices $u$ and $v$ in $G$ is the number of edges of a shortest path in $G$ connecting $u$ and $v$. The eccentricity $e_G(v)$ of a vertex $v$ is the maximum distance between

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and any other vertex, that is \( \max \{ \text{dist}_G(v, w) \mid w \in V \} \). The eccentric connectivity index \( \xi^c(G) \) of \( G \) is defined by

\[
\xi^c(G) = \sum_{v \in V} d_G(v) e_G(v).
\]

This index was introduced by Sharma et al. in [9] and successfully used for mathematical models of biological activities of diverse nature [2, 3, 5, 7, 8]. Recently, Hauweele et al. [4] have characterized those graphs which have the largest eccentric connectivity index among all connected graphs of given order \( n \). These results are summarized in Table 1, where

- \( K_n \) is the complete graph of order \( n \);
- \( P_n \) is the path of order \( n \);
- \( W_n \) is the wheel of order \( n \), i.e., the graph obtained by joining a vertex to all vertices of a cycle of order \( n - 1 \);
- \( M_n \) is the graph obtained from \( K_n \) by removing a maximum matching and, if \( n \) is odd, an additional edge adjacent to the unique vertex that still degree \( n - 1 \);
- \( E_{n,D} \) is the graph constructed from a path \( u_0 - u_1 - \ldots - u_D \) by joining each vertex of a clique \( K_{n-D-1} \) to \( u_0, u_1 \), and \( u_2 \).

### Table 1: Largest eccentric connectivity index for a fixed order \( n \)

<table>
<thead>
<tr>
<th>( n )</th>
<th>optimal graphs</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>( K_1 )</td>
</tr>
<tr>
<td>2</td>
<td>( K_2 )</td>
</tr>
<tr>
<td>3</td>
<td>( K_3 ) and ( P_3 )</td>
</tr>
<tr>
<td>4</td>
<td>( M_4 )</td>
</tr>
<tr>
<td>5</td>
<td>( M_5 ) and ( W_5 )</td>
</tr>
<tr>
<td>6</td>
<td>( M_6 )</td>
</tr>
<tr>
<td>7</td>
<td>( M_7 )</td>
</tr>
<tr>
<td>8</td>
<td>( M_8 ) and ( E_{8,4} )</td>
</tr>
<tr>
<td>( \geq 9 )</td>
<td>( E_{n,[\frac{n}{2}]}+1 )</td>
</tr>
</tbody>
</table>

In addition to the above-mentioned graphs, we will also consider the following ones:

- \( C_n \) is the chordless cycle of order \( n \);
- \( S_{n,x} \) is the graph of order \( n \) obtained by linking all vertices of a stable set of \( n - x \) vertices with all vertices of a clique \( K_x \). The graph \( S_{n,1} \) is called a star.

Also, for \( n \geq 4 \) and \( p \leq n - 3 \), let \( H_{n,p} \) be the graph of order \( n \) obtained by adding a dominating vertex (i.e., a vertex linked to all other vertices) to the graph or order \( n - 1 \) having \( p \) vertices of degree 0, and

- \( n - 1 - p \) vertices of degree 1 if \( n - p \) is odd;
- \( n - 2 - p \) vertices of degree 1 and one vertex of degree 2 if \( n - p \) is even.

For illustration, \( H_{8,3} \) and \( H_{9,3} \) are drawn on Figure 1. Note that \( H_{4,0} \simeq S_{4,2} \). Moreover, \( H_{4,0} \) has two dominating vertices while \( H_{4,1} \) and \( H_{n,p} \) have exactly one dominating vertex for all \( n \geq 5 \) and \( p \leq n - 3 \).
In this paper, we first give an alternative proof to a result of Zhou and Du [12] showing that the stars are the only graphs with smallest eccentric connectivity index among all connected graphs of given order $n \geq 4$. These graphs have $n - 1$ pending vertices (i.e., vertices of degree 1). We then consider all pairs $(n, p)$ of integers with $p \leq n - 1$ and characterize the graphs with smallest eccentric connectivity index among all connected graphs of order $n$ with $p$ pending vertices.

\section{Minimizing $\xi^c$ for graphs with fixed order}

$K_1$ and $K_2$ are the only connected graphs with 1 and 2 vertices, respectively, while $K_3$ and $P_3$ are the only connected graphs with 3 vertices. Since $\xi^c(K_3) = \xi^c(P_3) = 6$, all connected graphs of given order $n \leq 3$ have the same eccentric connectivity index. From now on, we therefore only consider connected graphs with fixed order $n \geq 4$. A proof of the following theorem was already given by Zhou and Du in [12]. Ours is slightly different.

**Theorem 1.** Let $G$ be a connected graph of order $n \geq 4$. Then $\xi^c(G) \geq 3(n - 1)$, with equality if and only if $G \simeq S_{n,1}$.

**Proof.** Let $x$ be the number of dominating vertices (i.e., vertices of degree $n - 1$) in $G$. We distinguish three cases.

- If $x = 1$, then let $u$ be the dominating vertex in $G$. Clearly, $e_G(u) = 1$ and $d_G(u) = n - 1$. All vertices $v \neq u$ have eccentricity $e_G(v) = 2$, while their degree is at least 1 (since $G$ is connected). Hence, $\xi^c(G) \geq (n - 1) + 2(n - 1) = 3(n - 1)$, with equality if and only if all $v \neq u$ have degree 1, i.e., $G \simeq S_{n,1}$.

- If $x > 1$, then all dominating vertices $u$ have $d_G(u)e_G(u) = n - 1$, while all non-dominating vertices $v$ have $d_G(v) \geq x \geq 2$ and $e_G(v) \geq 2$, which implies $d_G(u)e_G(u) \geq 4$. If $n = 4$, we therefore have $\xi^c(G) \geq 3n > 3(n - 1)$, while if $n > 4$, we have $\xi^c(G) \geq 2(n - 1) + 4(n - 2) = 6n - 10 > 3(n - 1)$.

- If $x = 0$, then every pending vertex $v$ has $e_G(v) \geq 3$ since its only neighbor is a non-dominating vertex. Since the eccentricity of the non-pending vertices is at least two, we have $d_G(v)e_G(v) \geq 3$ for all vertices $v$ in $G$, which implies $\xi^c(G) \geq 3n > 3(n - 1)$.

\hfill $\square$

Stars have $n - 1$ pending vertices. As will be shown in the next section, a similar result is more challenging when the total number of pending vertices is fixed to a value strictly smaller than $n - 2$. 

Figure 1: Two graphs with $p = 3$ pending vertices.
3 Minimizing $\xi^c$ for graphs with fixed order and fixed number of pending vertices

Let $G$ be a connected graph of order $n \geq 4$ with $p$ pending vertices. Clearly, $p \leq n - 1$, and $G \simeq S_{n,1}$ if $p = n - 1$. For $p = n - 2$, let $u$ and $v$ be the two non-pending vertices. Note that $u$ is adjacent to $v$ since $G$ is connected. Clearly, $G$ is obtained by linking $x \leq n - 3$ vertices of a stable set $S$ of $n - 2$ vertices to $u$, and the $n - 2 - x$ other vertices of $S$ to $v$. The $n - 2$ pending vertices $w$ have $d_G(w) = 1$ and $e_G(w) = 3$, while $e_G(u) = e_G(v) = 2$ and $d_G(u) + d_G(v) = n$. Hence $\xi^c(G) = 3(n - 2) + 2n = 5n - 6$ for all graphs of order $n$ with $n - 2$ pending vertices.

The above observations show that all graphs of order $n$ with a fixed number $p \geq n - 2$ of pending vertices have the same eccentric connectivity index. As will be shown, this is not the case when $n \geq 4$ and $p \leq n - 3$. We will prove that $H_{n,p}$ is almost always the unique graph minimizing the eccentric connectivity index. Note that

$$\xi^c(H_{n,p}) = \begin{cases} n - 1 + 2p + 4(n - p - 1) = 5n - 2p - 5 & \text{if } n - p \text{ is odd} \\ n - 1 + 2p + 4(n - p - 2) + 6 = 5n - 2p - 3 & \text{if } n - p \text{ is even.} \end{cases}$$

**Theorem 2.** Let $G$ be a connected graph of order $n \geq 4$ with $p \leq n - 3$ pending vertices and one dominating vertex. Then $\xi^c(G) \geq \xi^c(H_{n,p})$, with equality if and only if $G \simeq H_{n,p}$.

**Proof.** The dominating vertex $u$ in $G$ has $d_G(u)e_G(u) = n - 1$, the pending vertices $v$ have $d_G(v)e_G(v) = 2$, and the other vertices $w$ have $e_G(w) = 2$ and $d_G(w) \geq 2$. Hence, $\xi^c(G)$ is minimized if all non-pending and non-dominating vertices have degree 2, except one that has degree 3 if $n - p - 1$ is odd. In other words, $\xi^c(G)$ is minimized if and only if $G \simeq H_{n,p}$. \hfill $\square$

**Theorem 3.** Let $G$ be a connected graph of order $n \geq 4$, with at least two dominating vertices.

- If $n = 4$ then $\xi^c(G) \geq 12$, with equality if and only if $G \simeq K_4$.
- If $n = 5$ then $\xi^c(G) \geq 20$, with equality if and only if $G \simeq S_{5,2}$ or $G \simeq K_5$.
- If $n \geq 6$ then $\xi^c(G) \geq 6n - 10$, with equality if and only if $G \simeq S_{n,2}$.

**Proof.** Let $x$ be the number of dominating vertices in $G$. Then $d_G(u)e_G(u) = n - 1$ for all dominating vertices $u$, while $e_G(v) = 2$ and $d_G(v) \geq x$ for all other vertices $v$. Hence, $\xi^c(G) \geq -2x^2 + x(3n - 1)$.

- If $n = 4$ then $\xi^c(G) \geq f(x) = -2x^2 + 11x$. Since $2 \leq x \leq 4$, $f(2) = 14$, $f(3) = 15$, and $f(4) = 12$, we conclude that $\xi^c(G) \geq 12$, with equality if and only if $x = 4$, which is the case when $G \simeq K_4$.
- If $n = 5$ then $\xi^c(G) \geq f(x) = -2x^2 + 14x$. Since $2 \leq x \leq 5$, $f(2) = f(5) = 20$ and $f(3) = f(4) = 24$, we conclude that $\xi^c(G) \geq 20$, with equality if and only if $x = 2$ or 5, which is the case when $G \simeq S_{5,2}$ or $G \simeq K_5$.
- If $n \geq 6$ then $-2x^2 + x(3n - 1)$ is minimized for $x = 2$, which is the case when $G \simeq S_{n,2}$. \hfill $\square$

**Theorem 4.** Let $G$ be a connected graph of order $n \geq 4$, with $p \leq n - 3$ pending vertices and no dominating vertex. Then $\xi^c(G) > \xi^c(H_{n,p})$ unless $n = 5$, $p = 0$ and $G \simeq C_5$, in which case $\xi^c(G) = \xi^c(H_{n,0}) = 20$. 

4
Proof. Let $U$ be the subset of vertices $u$ in $G$ such that $d_G(u) = e_G(u) = 2$. If $U$ is empty, then all non-pending vertices $v$ in $G$ have $d_G(v) \geq 2$ and $e_G(v) \geq 2$ (since $G$ has no dominating vertex), and at least one of these two inequalities is strict, which implies $d_G(u)e_G(u) \geq 6$. Also, every pending vertex $w$ has $e_G(w) \geq 3$ since their only neighbor is not dominant. Hence, $\xi^c(G) \geq 6(n - p) + 3p = 6n - 3p$. Since $p \leq n - 3$, we have $\xi^c(G) \geq 5n - 2p + 3 > \xi^c(H_{n,p})$.

So, assume $U \neq \emptyset$. Let $u$ be a vertex in $U$, and let $v, w$ be its two neighbors. Also, let $A = N(v) \setminus (N(w) \cup \{w\})$, $B = (N(v) \cup N(w)) \setminus \{u\}$, and $C = N(w) \setminus (N(v) \cup \{v\})$. Since $e_G(u) = 2$, all vertices of $G$ belong to $A \cup B \cup C \cup \{u, v, w\}$. We finally define $B'$ as the subset of $B$ that contains all vertices $b$ of $B$ with $d_G(b) = 2$ (i.e., their only neighbors are $v$ and $w$).

Case 1: $v$ is adjacent to $w$.

$A \neq \emptyset$ else $w$ is a dominating vertex, and $C \neq \emptyset$ else $v$ is dominating. Let $G'$ be the graph obtained from $G$ by replacing every edge linking $v$ to a vertex $a \in A$ with an edge linking $w$ to $a$, and by removing all edges linking $v$ to a vertex of $B \setminus B'$. Clearly, $G'$ is also a connected graph of order $n$ with $p$ pending vertices, and $w$ is the only dominating vertex in $G'$. It follows from Theorem 2 that $\xi^c(G') \geq \xi^c(H_{n,p})$. Also,

- $d_G(u) = d_G'(u)$ and $e_G(u) = e_G'(u)$;
- $d_G(x) = d_G'(x)$ and $e_G(x) \geq e_G'(x)$ for all $x \in A \cup C$;
- $d_G(x) = d_G'(x)$ and $e_G(x) = e_G'(x)$ for all $x \in B'$;
- $d_G(x) > d_G'(x)$ and $e_G(x) = e_G'(x)$ for all $x \in B \setminus B'$.

Hence,

$$
\sum_{x \in A \cup B \cup C \cup \{u\}} d_G(x)e_G(x) \geq \sum_{x \in A \cup B \cup C \cup \{u\}} d_G'(x)e_G'(x).
$$

Moreover,

- $d_G(v)e_G(v) + d_G(w)e_G(w) = 2(|A| + |B| + 2) + 2(|C| + |B| + 2) = 2|A| + 2|B| + 2|C| + 8$;
- $d_G'(v)e_G'(v) + d_G'(w)e_G'(w) = 2(|B'| + 2) + |A| + |B| + |C| + 2$.

We therefore have

$$
\xi^c(G) - \xi^c(G') = \sum_{x \in A \cup B \cup C \cup \{u\}} d_G(x)e_G(x) + (d_G(v)e_G(v) + d_G(w)e_G(w))
$$

$$
- \sum_{x \in A \cup B \cup C \cup \{u\}} d_G'(x)e_G'(x) - (d_G'(v)e_G'(v) + d_G'(w)e_G'(w))
$$

$$
\geq (2|A| + 2|B| + 2|C| + 8) - (2(|B'| + 2) + |A| + |B| + |C| + 2)
$$

$$
= |A| + |C| + 3(|B'| + |B \setminus B'|) - 2|B'| + 2
$$

$$
= |A| + |C| + |B'| + 3|B \setminus B'| + 2 > 0.
$$

This implies $\xi^c(G) > \xi^c(G') \geq \xi^c(H_{n,p})$.

Case 2: $v$ is not adjacent to $w$, and both $A \cup (B \setminus B')$ and $C \cup (B \setminus B')$ are nonempty.

Let $G'$ be the graph obtained from $G$ by adding an edge linking $v$ to $w$, by replacing every edge linking $v$ to a vertex $a \in A$ with an edge linking $w$ to $a$, and by removing all edges linking $v$ to a vertex of $B \setminus B'$. Clearly, $G'$ is also a connected graph of order $n$ with $p$ pending vertices. As in the previous case, we have

$$
\sum_{x \in A \cup B \cup C \cup \{u\}} d_G(x)e_G(x) \geq \sum_{x \in A \cup B \cup C \cup \{u\}} d_G'(x)e_G'(x).
$$
Moreover, $e_G(v) \geq 2$ and $e_G(w) \geq 2$, while $e_{G'}(v) \leq 2$ and $e_{G'}(w) = 1$, which implies

- $d_G(v)e_G(v) + d_G(w)e_G(w) \geq 2(|A| + |B| + 1) + 2(|C| + |B| + 1) = 2|A| + 4|B| + 2|C| + 4$;
- $d_{G'}(v)e_{G'}(v) + d_{G'}(w)e_{G'}(w) \leq 2(|B'| + 2) + |A| + |B| + |C| + 2$.

We therefore have

$$\xi^c(G) - \xi^c(G') \geq (2|A| + 4|B| + 2|C| + 4) - (2(|B'| + 2) + |A| + |B| + |C| + 2)$$


If $B \setminus B' \neq \emptyset$, $w$ is the only dominating vertex in $G'$, and $\xi^c(G) - \xi^c(G') > 0$. It then follows from Theorem 2 that $\xi^c(G) > \xi^c(G') \geq \xi^c(H_{n,p})$. So assume $B \setminus B' = \emptyset$. Since $A \cup (B \setminus B') \neq \emptyset$, and $C \cup (B \setminus B') \neq \emptyset$, we have $A \neq \emptyset$ and $C \neq \emptyset$. Hence, once again, $w$ is the only dominating vertex in $G'$, and we know from Theorem 2 that $\xi^c(G') \geq \xi^c(H_{n,p})$.

- If $|B'| \geq 1$, $|A| \geq 2$ or $|C| \geq 2$, then $\xi^c(G) > \xi^c(G') \geq \xi^c(H_{n,p})$.
- If $|B'| = 0$ and $|A| = |C| = 1$, there are two possible cases:
  - if the vertex in $A$ is not adjacent to the vertex in $C$, then $n = 5$, $p = 2$, $G \simeq P_5$ and $G' \simeq H_{5,2}$. Hence, $\xi^c(G) = 24 > 16 = \xi^c(H_{n,p})$;
  - if the vertex in $A$ is adjacent to the vertex in $C$, then $n = 5$, $p = 0$, $G \simeq C_5$ and $G' \simeq H_{5,2}$. Hence, $\xi^c(G) = \xi^c(H_{n,p}) = 20$;

Case 3: $v$ is not adjacent to $w$, and at least one of $A \cup (B \setminus B')$ and $C \cup (B \setminus B')$ is empty. Without loss of generality, suppose $A \cup (B \setminus B') = \emptyset$. We distinguish two subcases.

Case 3.1: $B' = \emptyset$.

Since $n \geq 4$, $C \neq \emptyset$. Also, since $p \leq n - 3$, there is a non-pending vertex $r \in C$. Let $G'$ be the graph obtained from $G$ by removing the edge linking $u$ and $v$ and by linking $v$ to $u$ and to $r$. Note that $G'$ is a connected graph of order $n$ with $p$ pending vertices: while $v$ was pending in $G$, but not $u$, the situation is the opposite in $G'$. Note also that Theorem 2 implies $\xi^c(G') \geq \xi^c(H_{n,p})$ since $w$ is the only dominating vertex in $G'$. We then have:

- $d_G(u) = 2, d_{G'}(u) = 1$ and $e_G(u) = e_{G'}(u) = 2$, which gives $d_G(u)e_G(u) - d_{G'}(u)e_{G'}(u) = 2$;
- $d_G(v) = 1, d_{G'}(v) = 2$ and $e_G(v) = 2$, which gives $d_G(v)e_G(v) - d_{G'}(v)e_{G'}(v) = -1$;
- $d_G(w) = n - 2, d_{G'}(w) = n - 1$ and $e_G(w) = 2$ and $e_{G'}(w) = 1$, which gives $d_G(w)e_G(w) - d_{G'}(w)e_{G'}(w) = n - 3$;
- $d_G(r) = d_{G'}(r) + 1, e_G(r) = 3$ and $e_{G'}(w) = 2$, which gives $d_G(r)e_G(r) - d_{G'}(r)e_{G'}(r) = d_G(r) - 2$;
- $d_{G'}(c) = d_G(c)$ and $e_{G'}(c) > e_{G'}(c)$ for all $c \in (C \setminus \{r\})$. Since $r$ has a neighbor in $C$ of degree at least 2, we have $\sum_{c \in C \setminus \{r\}} (d_{G'}(c)e_{G'}(c) - d_G(c)e_G(c)) \geq 2$.

Hence, $\xi^c(G) - \xi^c(G') \geq 2 - 1 + n - 3 + d_G(r) - 2 + 2 > 0$, which implies $\xi^c(G) > \xi^c(G') \geq \xi^c(H_{n,p})$. 

\[ \text{6} \]
Case 3.2: \( B' = \emptyset \).
Let \( b_1, \ldots, b_{|B'|} \) be the vertices in \( B' \). Remember that the unique neighbors of these vertices are \( v \) and \( w \). Let \( G' \) be the graph obtained from \( G \) as follows. We first add an edge linking \( v \) to \( w \). Then, for every odd \( i < |B'| \), we add an edge linking \( b_i \) to \( b_{i+1} \) and remove the edges linking \( v \) to \( b_i \) and to \( b_{i+1} \). We then have

- \( d_G(x) = d_{G'}(x) \) and \( e_G(x) = e_{G'}(x) \) for all \( x \in B' \cup C \cup \{w\} \);
- \( d_G(v) = |B'| + 1, d_{G'}(v) \leq 3, e_G(v) \geq 2, \) and \( e_{G'}(v) \leq 2 \);
- \( d_G(w) = |B'| + |C| + 1, d_{G'}(w) = |B'| + |C| + 2, e_G(w) = 2, \) and \( e_{G'}(w) = 1 \).

Hence,

\[
\xi^c(G) - \xi^c(G') = d_G(v)e_G(v) + d_G(w)e_G(w) - d_{G'}(v)e_{G'}(v) - d_{G'}(w)e_{G'}(w) \\
\geq 2(|B'| + 1) + 2(|B'| + |C| + 1) - 6 - (|B'| + |C| + 2) \\
= 3|B'| + |C| - 4.
\]

If \( |B'| \geq 2 \) or \( |C| \geq 2 \), then \( \xi^c(G) - \xi^c(G') > 0 \), and since \( w \) is then the only dominating vertex in \( G' \), we know from Theorem 2 that \( \xi^c(G) > \xi^c(G') \geq \xi^c(H_{n,p}) \). So, assume \( |B'| = 1 \) and \( |C| \leq 1 \):

- if \( |C| = 0 \) then \( n = 4, p = 0, G \cong C_4 \) and \( G' \cong H_{4,0} \) which implies \( \xi^c(G) = 16 > 14 = \xi^c(H_{n,p}) \);
- if \( |C| = 1 \) then \( n = 5, p = 1, \xi^c(G) = 23 \) and \( G' \cong H_{5,1} \) which implies \( \xi^c(G) > 20 = \xi^c(H_{n,p}) \).

\[ \square \]

We can now combine these results as follows. Assume \( G \) is a connected graph of order \( n \) with \( p \) pending vertices. If \( p \geq 1 \), then \( G \) has at most one dominating vertex, and it follows from Theorems 2 and 3 that \( H_{n,p} \) is the only graph with maximum eccentric connectivity index. If \( p = 0 \) and \( n = 4 \), then \( G \) cannot contain exactly one dominating vertex, and Theorems 3 and 4 show that \( K_4 \) is the only graph with maximum eccentric connectivity index. If \( p = 0 \) and \( n = 5 \), Theorems 2, 3 and 4 show that \( H_{5,0}, S_{5,2}, K_5 \) and \( C_5 \) are the only candidates to minimize the eccentric connectivity index, and since \( \xi^c(H_{5,0}) = \xi^c(S_{5,2}) = \xi^c(K_5) = \xi^c(C_5) = 20 \), the four graphs are the optimal ones. If \( p = 0 \) and \( n \geq 6 \) then we know from Theorems 3 and 3 that \( S_{n,2} \) and \( H_{n,0} \) are the only candidates to minimize the eccentric connectivity index. Since \( \xi^c(S_{n,2}) = 26 < 27 = \xi^c(H_{6,0}) \), \( \xi^c(S_{7,2}) = 32 > 30 = \xi^c(H_{7,0}) \) and \( \xi^c(S_{n,2}) = 6n - 10 > 5n - 3 \geq \xi^c(H_{n,0}) \) for \( n \geq 8 \), we deduce that \( S_{n,2} \) is the only graph with maximum eccentric connectivity index when \( n = 6 \) and \( p = 0 \), while \( H_{n,0} \) is the only optimal graph when \( n \geq 7 \) and \( p = 0 \). This is summarized in the following Corollary.

**Corollary 5.** Let \( G \) be a connected graph of order \( n \geq 4 \) with \( p \leq n - 3 \) pending vertices.

- If \( p \geq 1 \) then \( \xi^c(G) \geq \xi^c(H_{n,p}) \) with equality if and only if \( G \cong H_{n,p} \);
- If \( p = 0 \) then
  - if \( n = 4 \) then \( \xi^c(G) \geq 12 \), with equality if and only if \( G \cong K_4 \);
  - if \( n = 5 \) then \( \xi^c(G) \geq 20 \), with equality if and only if \( G \cong H_{5,0}, S_{5,2}, K_5 \) or \( C_5 \);
  - if \( n = 6 \) then \( \xi^c(G) \geq 26 \), with equality if and only if \( G \cong S_{6,2} \);
  - if \( n \geq 7 \) then \( \xi^c(G) \geq \xi^c(H_{n,0}) \), with equality if and only if \( G \cong H_{n,0} \).
4 Conclusion

We have characterized the graphs with smallest eccentric connectivity index among those of fixed order \( n \) and fixed or non-fixed number of pending vertices. Such a characterization for graphs with a fixed order \( n \) and a fixed size \( m \) was given in [12]. It reads as follows.

**Theorem 6.** Let \( G \) be a connected graph of order \( n \) with \( m \) edges, where \( n - 1 \leq m < \binom{n}{2} \). Also, let

\[
k = \left\lfloor \frac{2n - 1 - \sqrt{(2n - 1)^2 - 8m}}{2} \right\rfloor.
\]

Then \( \xi_c(G) \geq 4m - k(n - 1) \), with equality if and only if \( G \) has \( k \) dominating vertices and \( n - k \) vertices of eccentricity 2.

It is, however, an open question to characterize the graphs with largest eccentric connectivity index among those of fixed order \( n \) and fixed size \( m \). The following conjecture appears in [4], where \( E_{n,D,k} \) is the graph of order \( n \) constructed from a path \( u_0 - u_1 - \ldots - u_D \) by joining each vertex of a clique \( K_{n-D-1} \) to \( u_0 \) and \( u_1 \), and \( k \) vertices of the clique to \( u_2 \).

**Conjecture 7.** Let \( G \) be a connected graph of order \( n \) with \( m \) edges, where \( n - 1 \leq m \leq \binom{n-1}{2} \). Also, let

\[
D = \left\lfloor \frac{2n + 1 - \sqrt{17 + 8(m - n)}}{2} \right\rfloor \quad \text{and} \quad k = m - \left( \frac{n - D + 1}{2} \right) - D + 1.
\]

Then \( \xi_c(G) \leq \xi_c(E_{n,D,k}) \), with equality if and only if \( G \simeq E_{n,D,k} \) or \( D = 3, k = n - 4 \) and \( G \) is the graph constructed from a path \( u_0 - u_1 - u_2 - u_3 \), by joining \( 1 \leq i \leq n - 3 \) vertices of a clique \( K_{n-4} \) to \( u_0, u_1, u_2 \) and the \( n - 4 - i \) other vertices of \( K_{n-4} \) to \( u_1, u_2, u_3 \).

References


