

PHOEG Helps Obtaining Extremal Graphs

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Joint work with Gauvain Devillez and Hadrien Mélot

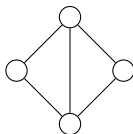
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CSD8, Mons, August 23, 2017

Introduction

We consider **simple undirected** graphs.



For a graph $G = (V, E)$,

- its **order** $|V|$ is denoted by n ;
- its **size** $|E|$ is denoted by m .

A **graph invariant** is a function on graphs that is constant on isomorphism classes.

Examples: order n , size m , chromatic number χ , maximum degree Δ , diameter D , planarity, ...

Extremal Graph Theory

Extremal Graph Theory aims to find bounds on a graph invariant under some constraints.

Generally, those constraints are of two types:

- restricting class of graphs (e.g., connected graphs, trees);
- fixing (and restricting) values of other invariants (e.g., size, maximum degree).

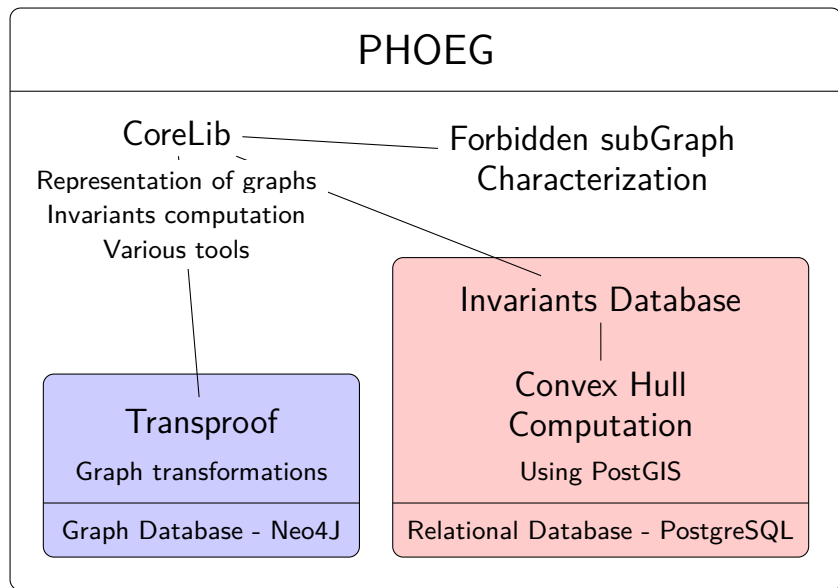
Results in Extremal Graph Theory mainly consists in

- giving bounds;
- characterizing graphs achieving these bounds.

Computer-assisted discovery

- **Context:** Computer-assisted Discovery in Extremal Graph Theory
- **Several existing systems:** Graph, Graffiti, AutoGraphiX, GraPHedron, . . .
 - exploit different ideas to help graph theorists
- **Objectives of this talk:**
 - presentation of PHOEG, a successor of GraPHedron
 - use of an illustrative problem (Eccentric Connectivity Index, ECI)
- **Remark:** work in progress
 - PHOEG is currently a prototype
 - the problem about ECI is not fully solved

Overview of PHOEG

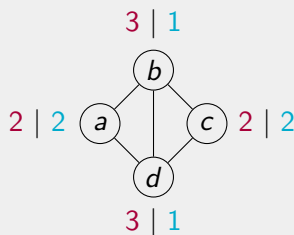


Eccentricity Connectivity Index

Let v be a vertex of a graph G , recall that:

- **degree** $d(v)$ = number of adjacent vertices of v ;
- **eccentricity** $\epsilon(v)$ = maximal distance between v and any other vertex.

Example



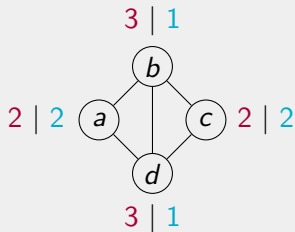
Eccentric Connectivity Index

Definition

The **Eccentric Connectivity Index** (ECI) of a graph G , denoted by $\xi^c(G)$, is

$$\xi^c(G) = \sum_{v \in V} d(v)\epsilon(v).$$

Example



$$\xi^c(G) = (2 \times 2 + 3 \times 1) \times 2 = 14$$

Eccentric Connectivity Index

History and motivation

- Sharma, Goswani and Madan introduced ξ^c in 1997 in Chemistry;
- Useful as a discriminating topological descriptor for Structure Properties and Structure Activity studies;
- Since 1997, more than 200 chemical papers about ξ^c : applications in drug design, prediction of anti-HIV activities, etc.
- However, the first mathematical paper with extremal properties on ξ^c was published only in 2010;
- Since 2010, about a dozen papers containing bounds on ξ^c .

Some Extremal Theory problem about ξ^c

Now, let's make extremal graph theory about ξ^c with the help of a computer.

First step: define a problem by choosing constraints.

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Several papers containing bounds on ξ^c — using various invariants as constraints — have been published (since 2010).

Some Extremal Theory problem about ξ^c

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Problem

Among connected graphs of order n and size m , what is the maximum possible value for ξ^c ?

Upper bound on ξ^c for connected graphs with fixed size

We define $E_{n,m}$ as follows :

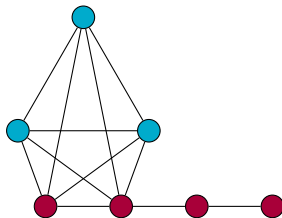
$$n = 7, m = 14$$

Upper bound on ξ^c for connected graphs with fixed size

We define $E_{n,m}$ as follows :

- The biggest possible clique without disconnecting the graph, leaving a path with the remaining vertices.

$$n = 7, m = 14$$

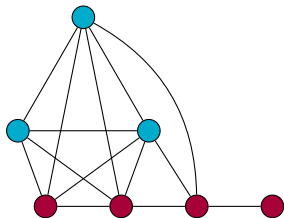


Upper bound on ξ^c for connected graphs with fixed size

We define $E_{n,m}$ as follows :

- The biggest possible clique without disconnecting the graph, leaving a path with the remaining vertices.
- Add remaining edges between vertices of the clique and the first vertex of the path.

$$n = 7, m = 14$$

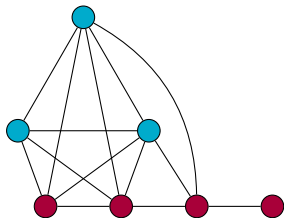


Upper bound on ξ^c for connected graphs with fixed size

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- The biggest possible clique without disconnecting the graph, leaving a path with the remaining vertices.
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$$n = 7, m = 14$$



This graph is unique for given n and m . We define $d_{n,m}$ as the diameter of $E_{n,m}$.

Conjecture of Zhang, Liu and Zhou

Conjecture (Zhang, Liu and Zhou, 2014)

Let G be a graph of order n and size m such that $d_{n,m} \geq 3$. Then,

$$\xi^c(G) \leq \xi^c(E_{n,m}),$$

with equality if and only if $G \simeq E_{n,m}$.

- The authors prove that the conjecture is true when $m = n - 1, n, \dots, n + 4$ (if n is large enough).
- There exists a “proof” published in a journal of University of Isfahan (Iran, 2014) but that is obviously wrong.

Conjecture of Zhang, Liu and Zhou

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with equality if and only if $G \simeq E_{n,m}$.

- Is the conjecture true?
- If yes, how to prove it?
- If no, how to improve or correct it?
- What about graphs such that $d_{n,m} < 3$?

How can the computer help?

In the following, we will show how PHOEG can help to study all of the preceding questions and to raise new ones.

P H_{elps} O_{btaining} E_{xtremal} G_{raphs}

PHOEG — the database part

- Former system (GraPHedron): graphs and invariant's values written sequentially in files;
- PHOEG uses a [PostgreSQL DB](#) with tens of millions of non-isomorphic graphs and invariants' values;
- Invariant's values are computed once (useful for NP-hard invariants);

Database of the invariants

- Each graph has its unique **signature** used as primary key (canonical form, thanks to Nauty by Brendan McKay), $sig(C_5) = "DqK"$, $sig(K_3) = "Bw"$.
- 12 millions simple graphs up to order 10, 8 millions cubic graphs up to order 22.

Graphs
signature
A_
A?
B?
BG
Bw
BW
C'
C^
C~
C?
C@

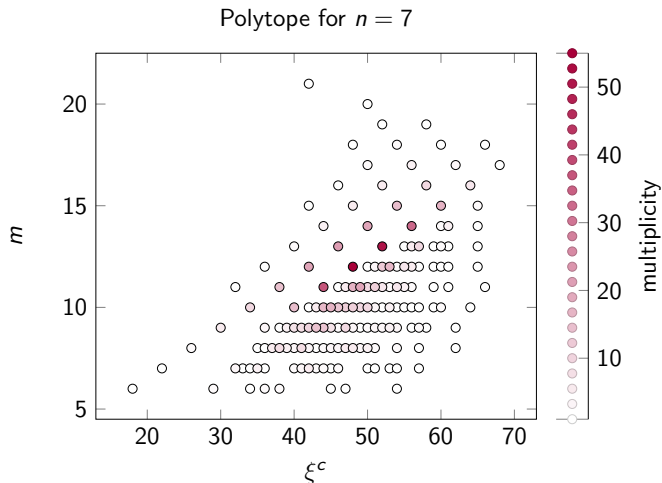
NumVertices	
signature	val
A_	2
A?	2
B?	3
BG	3
Bw	3
BW	3
C'	4
C^	4
C~	4
C?	4
C@	4

NumEdges	
signature	val
A_	1
A?	0
B?	0
BG	1
Bw	3
BW	2
C'	2
C^	5
C~	6
C?	0
C@	1

ECI	
signature	val
A_	2
BW	6
Bw	6
C^	14
C~	12
CF	9
CN	13
Cr	16
CR	14
D' [25
D' {	20

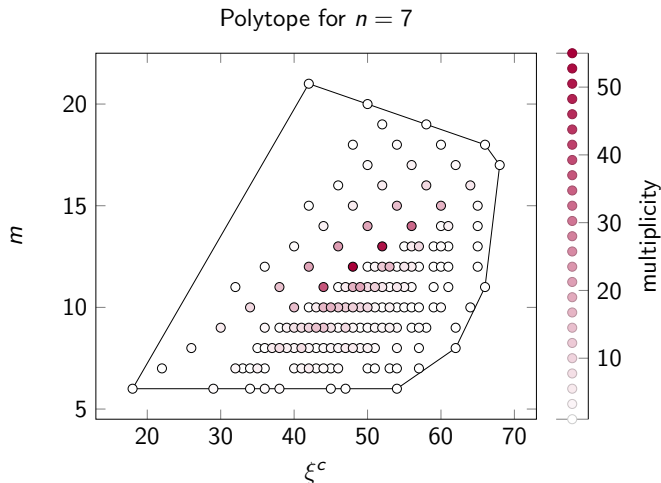
GraPHedron's main principle

- view graphs as points in the space of invariants;



GraPHedron's main principle

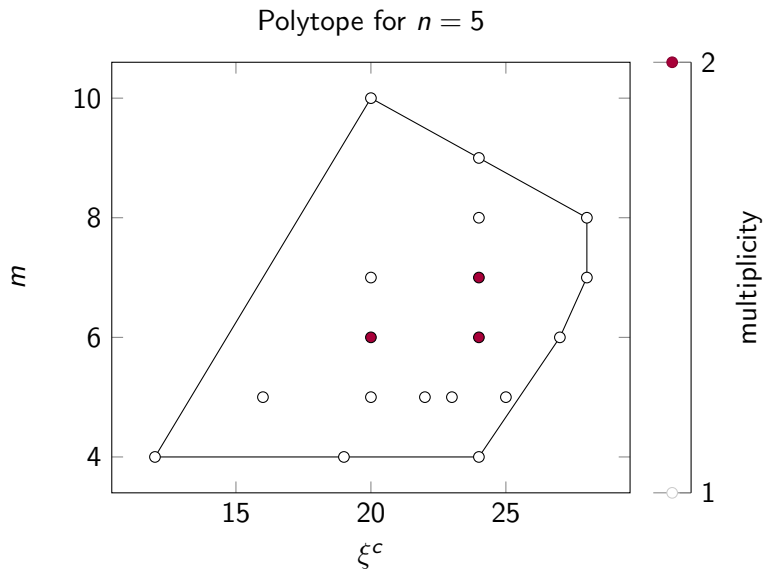
- view graphs as points in the space of invariants;
- compute the convex hull of these points (for small values of n).



Database query – Points, multiplicities and polytope

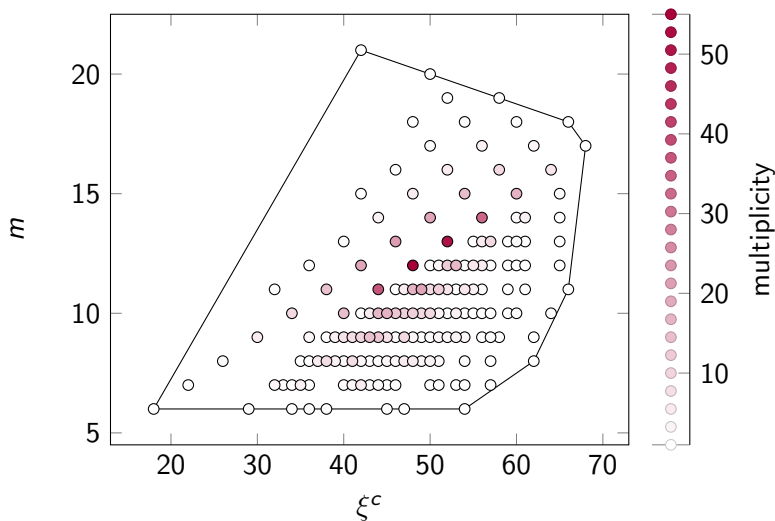
	eci	m	mult
	-----	-----	-----
SELECT P.val AS eci, num_edges.val AS m, COUNT(*) AS mult	47	8	5
FROM eci P	46	8	3
JOIN num_vertices USING(signature)	40	8	3
JOIN num_edges USING(signature)	32	7	3
WHERE num_vertices.val = 7	48	12	55
GROUP BY m, eci;	48	18	1
	61	14	4
	59	13	1
	48	11	17
SELECT ST_AsText(ST_ConvexHull(ST_Collect(ST_Point(eci, m))))	43	9	14
FROM poly;	47	6	1
	64	10	1
	59	11	1
st_astext	45	9	7
-----	38	6	2
POLYGON((18 6,42 21,66 18,68 17,66 11,62 8,54 6,18 6))	[...]		

Exploring ξ^c with PHOEG: polytopes



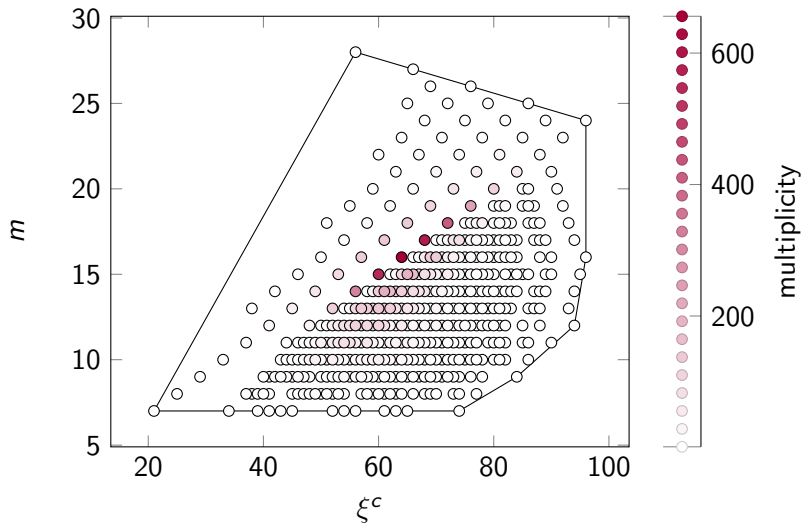
Exploring ξ^c with PHOEG: polytopes

Polytope for $n = 7$

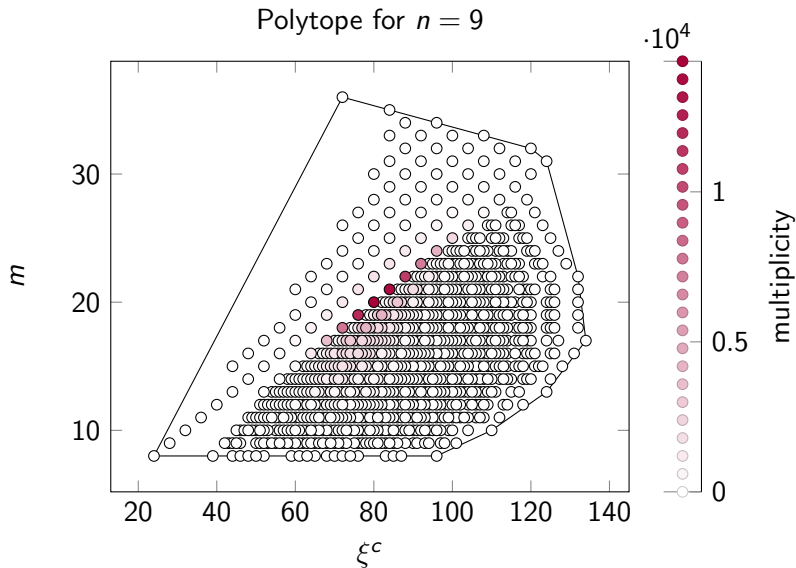


Exploring ξ^c with PHOEG: polytopes

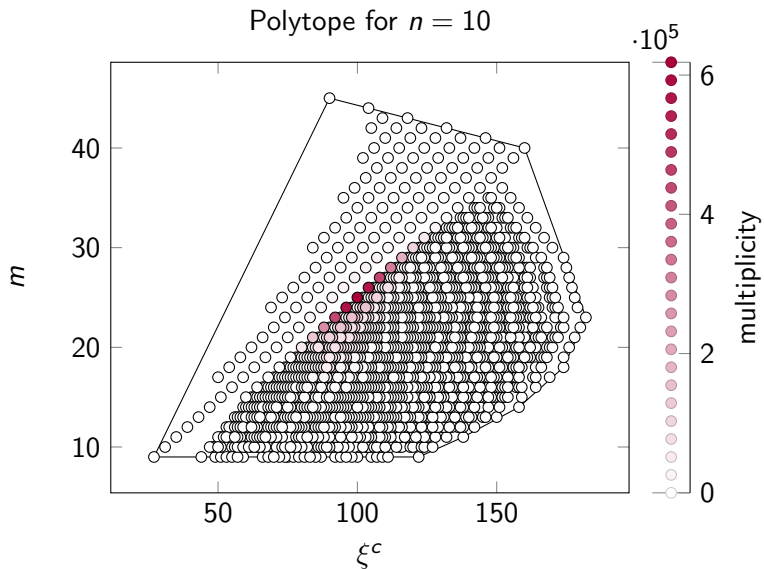
Polytope for $n = 8$



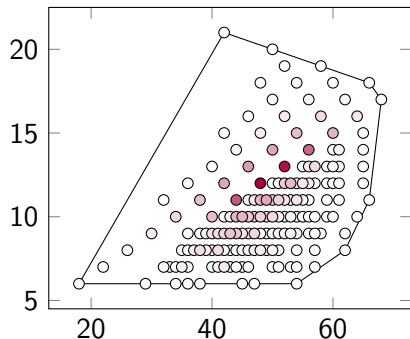
Exploring ξ^c with PHOEG: polytopes



Exploring ξ^c with PHOEG: polytopes

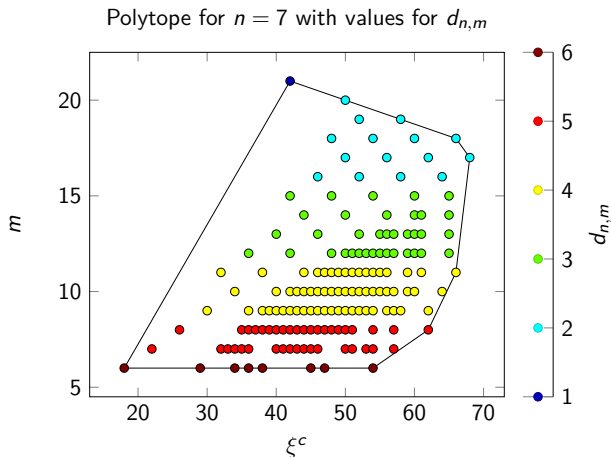


Observations and questions



- How to explain the grid?
- Is the conjecture of Zhang, Liu and Zhou true when $d_{n,m} \geq 3$?
- Upper bound when $d_{n,m} < 3$?

Coloring points with values of $d_{n,m}$



Recall that the conjecture is stated for $d_{n,m} \geq 3$. Is it true for $n = 7$?

Database query – Extremal graphs

```
WITH tmp AS (  
  SELECT n.val AS n, m.val AS m,  
         P.signature, P.val AS eci, d.val AS d,  
         rank() OVER (  
           PARTITION BY n.val, m.val  
           ORDER BY P.val DESC  
         ) AS pos  
  FROM num_vertices n  
  JOIN num_edges m USING(signature)  
  JOIN d USING(signature)  
  JOIN eci P USING(signature)  
  WHERE n.val = 7  
)  
SELECT signature AS sig, n, m, eci, d  
FROM tmp  
WHERE pos = 1 AND d >= 3  
ORDER BY n, m, d, eci;
```

sig	n	m	eci	d
F@IQO	7	6	54	6
F@'J_	7	7	57	5
FgCXW	7	8	62	5
FWCYw	7	9	62	4
FgCxw	7	10	64	4
F'Kyw	7	11	66	4
F'Kzw	7	12	65	3
F'Lzw	7	13	65	3
F'\zw	7	14	65	3
FJ] w	7	15	65	3
FJ\ w	7	15	65	3

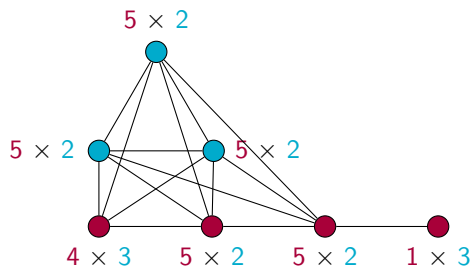
Database query – Extremal graphs

```
WITH tmp AS (  
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         P.signature, P.val AS eci, d.val AS d,  
         rank() OVER (  
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           ORDER BY P.val DESC  
         ) AS pos  
  FROM num_vertices n  
  JOIN num_edges m USING(signature)  
  JOIN d USING(signature)  
  JOIN eci P USING(signature)  
  WHERE n.val = 7  
)  
SELECT signature AS sig, n, m, eci, d  
FROM tmp  
WHERE pos = 1 AND d >= 3  
ORDER BY n, m, d, eci;
```

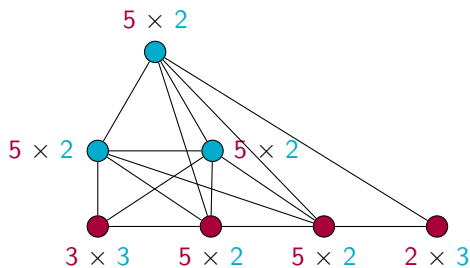
sig	n	m	eci	d
F@IQO	7	6	54	6
F@'J_	7	7	57	5
FgCXW	7	8	62	5
FWCYw	7	9	62	4
FgCxw	7	10	64	4
F'Kyw	7	11	66	4
F'Kzw	7	12	65	3
F'Lzw	7	13	65	3
F'\zw	7	14	65	3
FJ] w	7	15	65	3
FJ\ w	7	15	65	3

⇒ counter-example to the conjecture !
Extremal graphs are not always unique

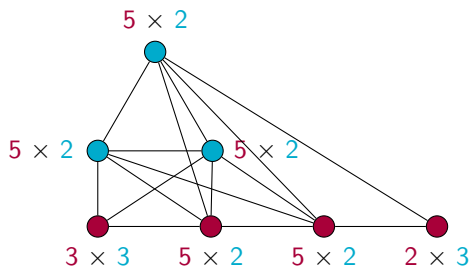
Counter-example ($n = 7$ and $m = 15$)



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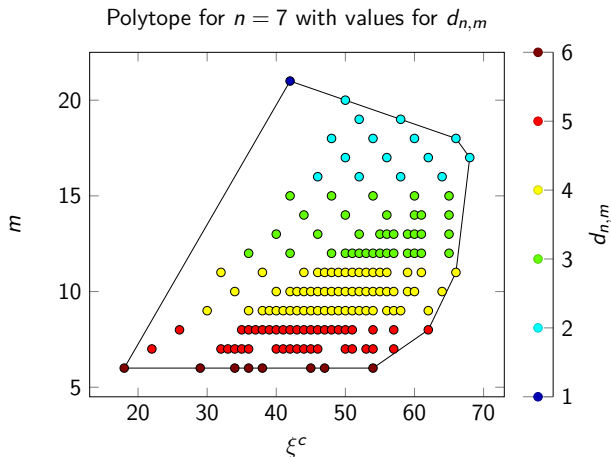


Counter-example ($n = 7$ and $m = 15$)



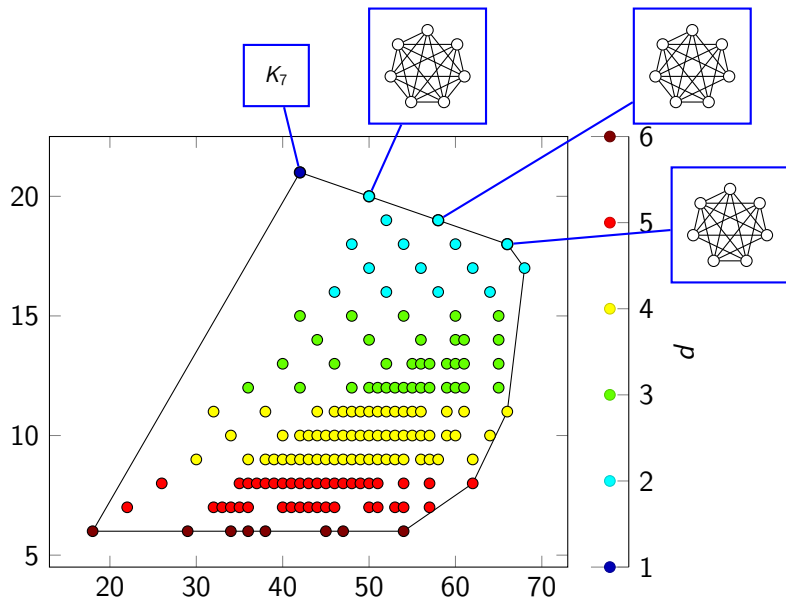
It is possible to construct counter-examples for any values of $n \geq 6$ (with $d_{n,m} = 3$).

Coloring points with values of $d_{n,m}$

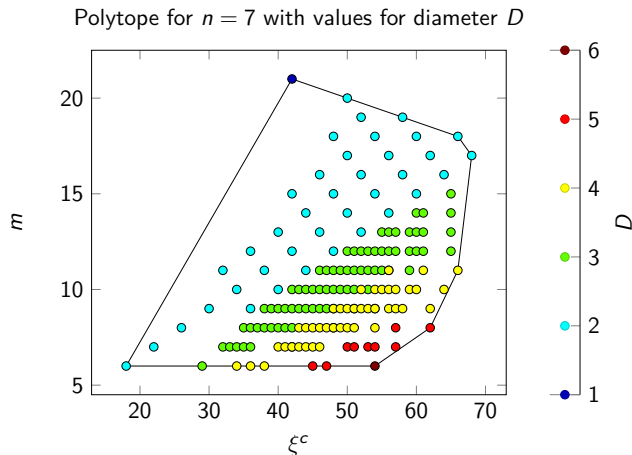


Upper bound when $d_{n,m} < 3$?

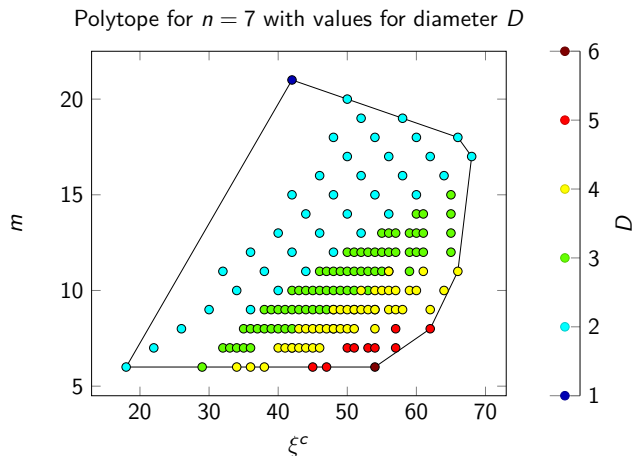
Upper facet of the polytope ($n = 7$)



Coloring points with values of the diameter



Coloring points with values of the diameter



Can the diameter explain the blue grid? Actually, yes!

A new tight upper bound when $d_{n,m} < 3$

Theorem

Let G be a graph of order n and size m . Then,

$$\xi^c(G) \leq n(n-1)(n-2) - 2m(n-3),$$

with equality if and only if G is the complement of a matching.

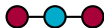
Note that the bound is valid for all graphs but can be tight only if

$$m \geq \binom{n}{2} - \left\lfloor \frac{n}{2} \right\rfloor,$$

(and thus $d_{n,m} < 3$).

Number of non-equivalent colorings

We note $P(G, k)$ the number of **non-equivalent** colorings of G that use **exactly** k colors.



$$P(P_3, 2) = 1$$



$$P(P_3, 3) = 1$$

Total number of non-equivalent colorings

Definition

The total number of non-equivalent colorings $\mathcal{P}(G)$ of a graph G is

$$\mathcal{P}(G) = \sum_{k=0}^n P(G, k) = \sum_{k=\chi(G)}^n P(G, k),$$

where $\chi(G)$ is the chromatic number of G .

Example: $\mathcal{P}(P_3) = P(P_3, 2) + P(P_3, 3) = 1 + 1 = 2$

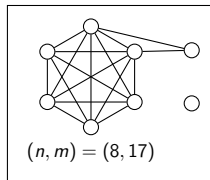
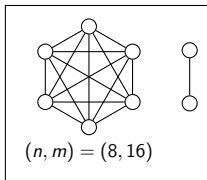
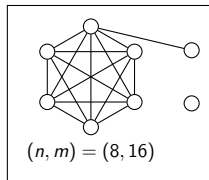
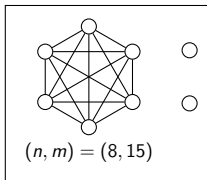
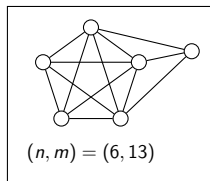
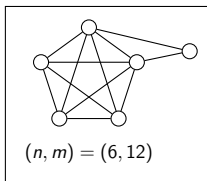
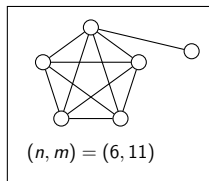
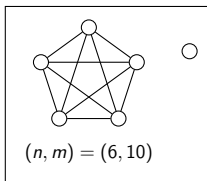
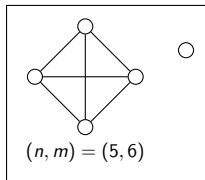
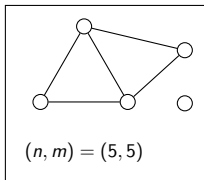
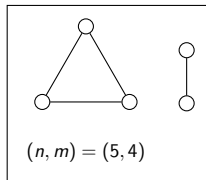
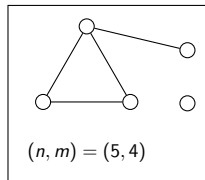
$\mathcal{P}(G)$ is the value of the σ -polynomial when $x = 1$ and is also known as the **Bell number** of a graph [Duncan & Peele, 2009].

The Min-NumCol-NumEdges Problem

Problem

What is minimum possible value of \mathcal{P} for graphs of fixed order n and size m and what are the graphs attaining those bounds ?

Some extremal graphs

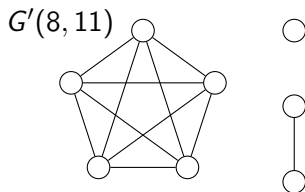
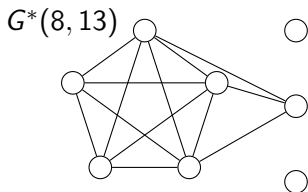


The extremal(?) graphs

Given n the order and m the size of graphs. Let t_k be the biggest triangular number such that $t_k \leq m$. We call $r_m = m - t_k$ the **remainder**.

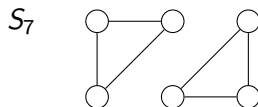
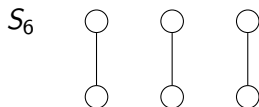
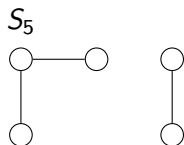
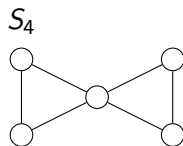
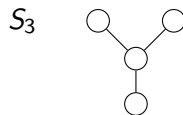
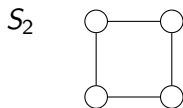
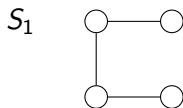
We define $G^*(n, m)$ as the unique graph formed from $K_{k+1} \cup \overline{K}_{n-k-1}$, where one (if any) vertex of \overline{K}_{n-k-1} is connected to r_m vertices of the clique.

If $r_m = 1$, and $n - k - 1 \geq 2$, we define $G'(n, m)$ as $K_{k+1} \cup \overline{K}_{n-k-1}$, where two vertices of $K_{k+1} \cup \overline{K}_{n-k-1}$ are connected.



Forbidden Graph Characterization

In this tool, we want a necessary *and* sufficient characterization of our graphs.



Concluding remarks

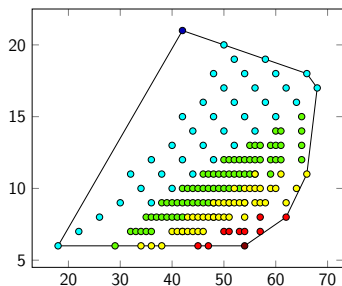
- Not only extremal graphs are useful to study extremal properties of an invariant
- Exact approach limited to small graphs ($n \leq 10$)
- However, dealing with small graphs has already shown to be very useful and led to numerous results (AutoGraphiX, GraPHedron)

Perspectives

- Invariants' DB allows a form of dynamic programming;
- Create a simple interface for queries, define a domain specific language;
- Allow easy visualization and manipulation of outputs (GUI, PDF, etc.);
- Go up in the order of graphs, relaxing the *exact* constraint.

Appendix

Understanding the grid of blue points



- Suppose $D(G) = 2$ (light blue points)
- For each vertex v , since $D(G) = 2$, either $\epsilon(v) = 1$ or $\epsilon(v) = 2$
- If $\epsilon(v) = 1$, then v is dominant and $d(v) = n - 1$
- Let k be the number of dominant vertices of G
- The sum of degrees of non dominant vertices is

$$2m - k(n - 1)$$

Thus,

$$\xi^c(G) = k(n - 1) + 2(2m - k(n - 1)) = 4m - k(n - 1),$$

that is maximum if $k = 0$ and, moreover, explain the grid.

Upper bound on ξ^c for connected graphs with fixed size

Definition

For positive integers n and m with $n - 1 \leq m \leq \binom{n}{2}$, let

$$d_{n,m} = \left\lfloor \frac{2n + 1 - \sqrt{17 + 8(m - n)}}{2} \right\rfloor.$$

In the following, we simply use d for $d_{n,m}$.

Definition

Let $E_{n,m}$ be the graph obtained from a clique K_{n-d-1} and a path $P_{d+1} = v_0 v_1 \dots v_d$ by joining each vertex of the clique to both v_d and v_{d-1} , and by joining

$$m - n + 1 - \binom{n-d}{2}$$

vertices of the clique to v_{d-2} .

Upper bound on ξ^c for connected graphs with fixed size

Example ($n = 5$)

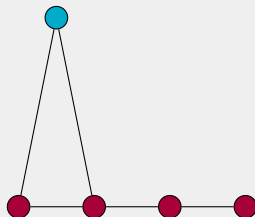
m	4	5	6	7	8	9	10
d	4	3	3	2	2	2	1
$n - d - 1$	0	1	1	2	2	2	3
# edges to v_{d-2}	0	0	1	0	1	2	0



Upper bound on ξ^c for connected graphs with fixed size

Example ($n = 5$)

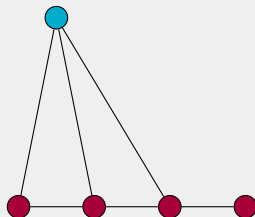
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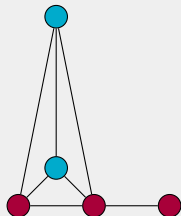
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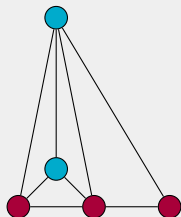
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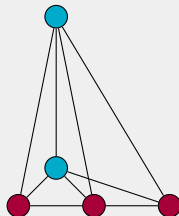
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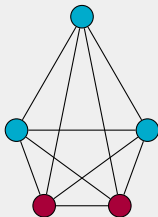
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What about other classes of graphs ?

Let's try to maximize ξ^c on cubic (3-regular) graphs.

```
SELECT t.n, t.signature, t.eci
FROM (
  SELECT n.val AS n, eci.signature, eci.val as eci,
         DENSE_RANK() OVER (
           PARTITION BY n.val
           ORDER BY eci.val DESC
         ) AS pos
  FROM cubic
  JOIN num_vertices n USING(signature)
  JOIN eccentric_connectivity_index eci USING(signature)
) t
WHERE t.pos = 1
ORDER BY t.n;
```

Maximize ξ^c on cubic graphs

n	signature	eci
4	C~	12
6	Es\o	36
6	E{Sw	36
8	Gv?IXW	72
8	Gs@ipo	72
10	Iv?GOKFY?	126
12	Kt?GOKFOAOeA	177
14	Mt?GO?@@_KgKOWM??	270
16	Ot?G?CA?WB'oO?O?b_@?E	348
18	Qv??W[K?G??@?B?B?A??'G?p??o	474
20	Sv?GW?@?W??@?B????G?J??w?w?M?BO??	573
22	Uv?G?CK?oE@_?H?E??G?C??C??W?@??@C_?KO??o	726

